

Scalable Mechanism Design for Multi-Agent Path Finding

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Talk at IJCAI '24, 09 August 2024



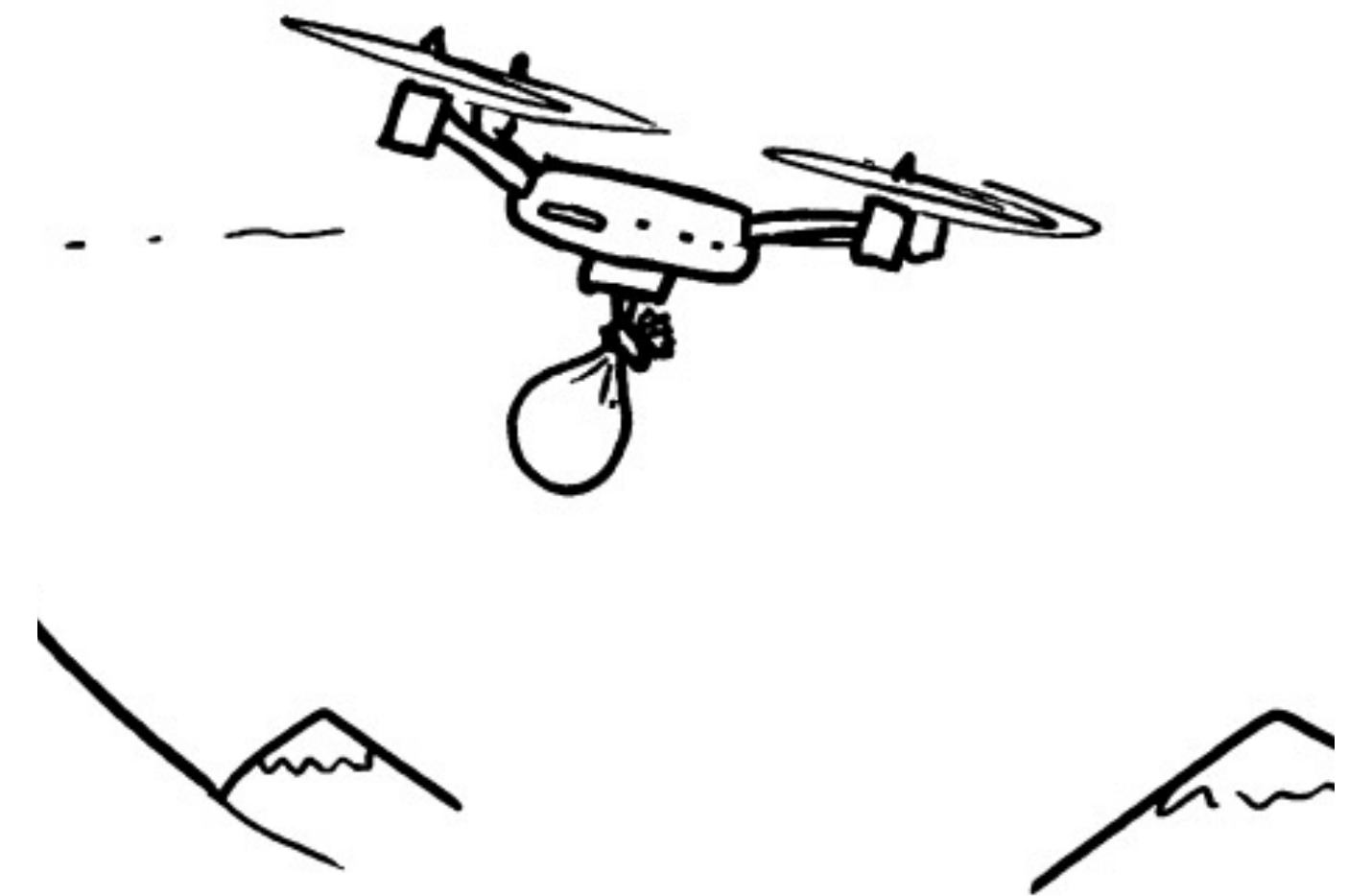
Competing for paths: a new domain

Combines

- **large scale**
- **mechanism design:**
self-interested agents
- **multi-agent path finding (MAPF):**
to calculate allocation

Traffic management: allocating

- road capacity to cars
- urban airspace to UAVs



Not your usual MAPF setting

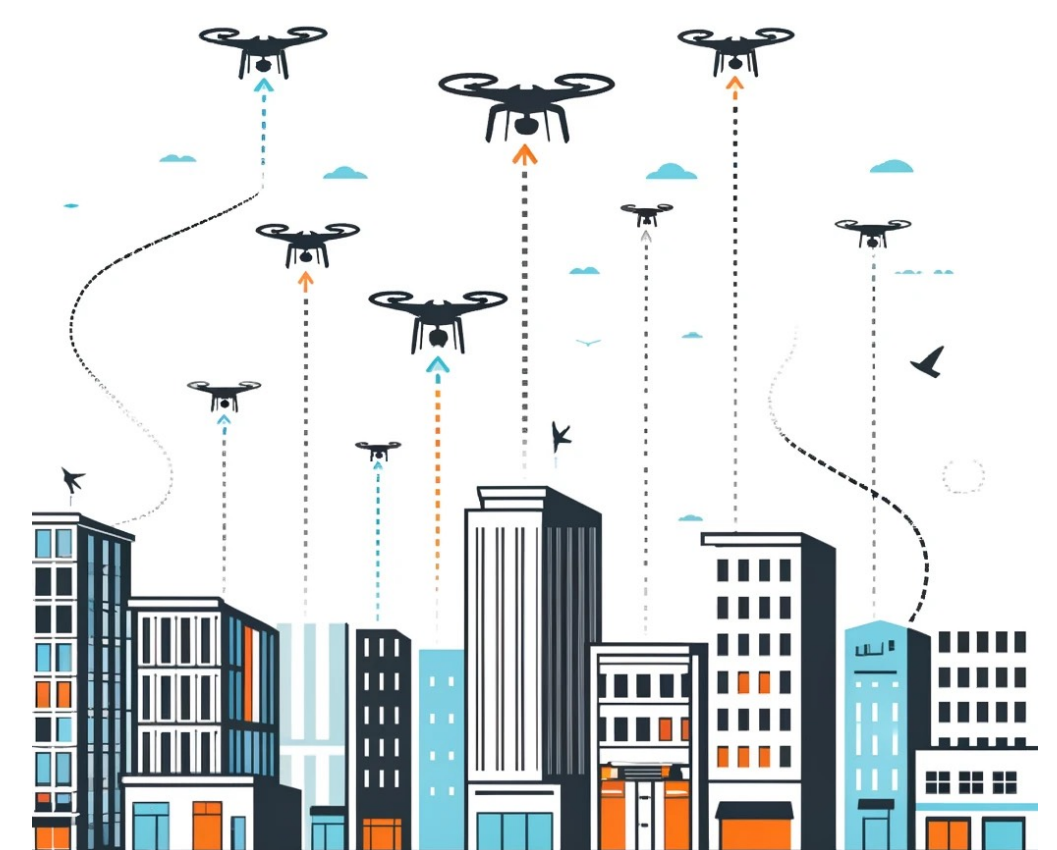
The usual MAPF setting is **cooperative**:

- A central allocator system knows all true agent's characteristics

Our setting is **non-cooperative**:

- The system queries agents for their characteristics
- Agents will lie if they can secure a faster route

Without incentive awareness: **worse overall solutions!**



Problem summary

Design mechanisms for allocating collision-free paths that are:

- **strategyproof:** incentivize agents to be honest
- **efficient:** find allocations with high consumer welfare
- **scalable**

Up until now: "pick 2"...

... we can do all three!

Why don't we just use auctions?

Mechanism designers love auctions, especially **the VCG auction**.

Ask agents how
much they value
the resource



Find the allocation
that's optimal for
these reports



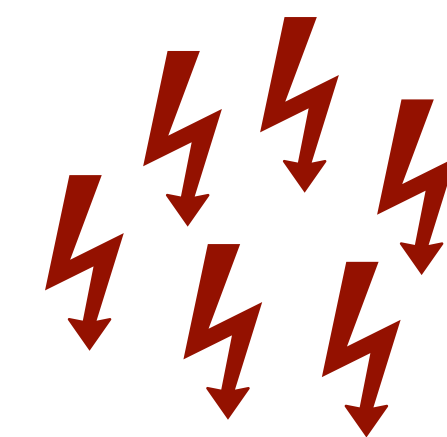
Charge clever
payments that make
lying unprofitable



**Many possible paths /
allocations / vertex
combinations to ask**



**Optimal MAPF is
hard!**



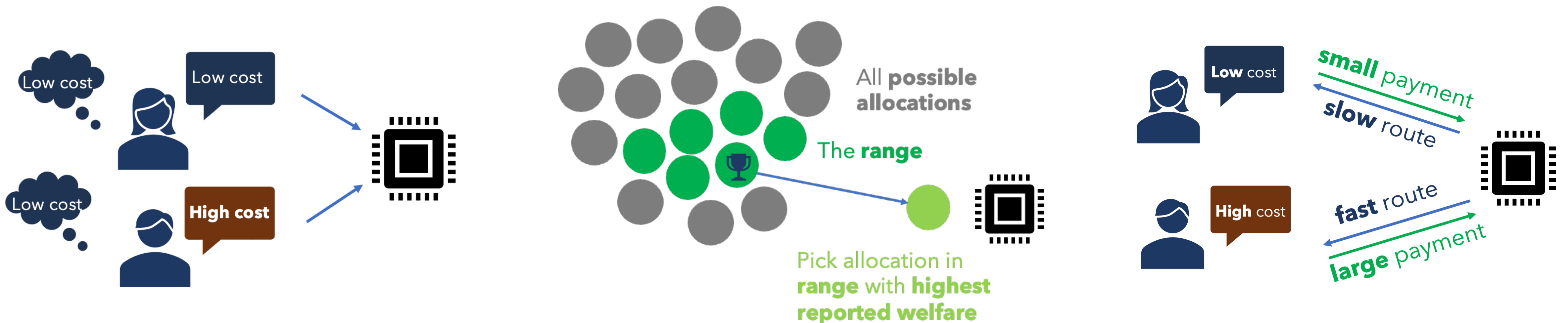
**Optimal MAPF is
hard! ($\times N$)**

The solution: maximal-in-range (MIR)

Ask agents for their start & goal, value for arrival, cost for delay

Suboptimal MAPF algo finds allocation that is **optimal within a range**

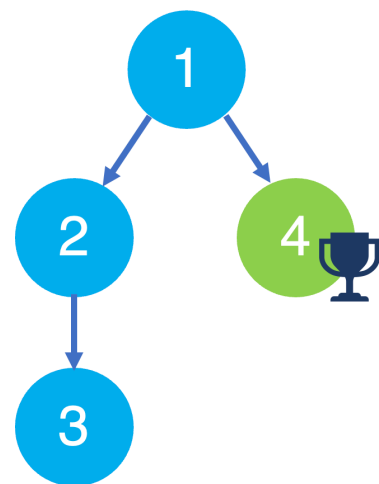
Charge VCG payments, with a twist that makes them "free"



The mechanism is **strategy-proof** and **individually rational**, as long as agents **cannot influence the range by misreporting!**

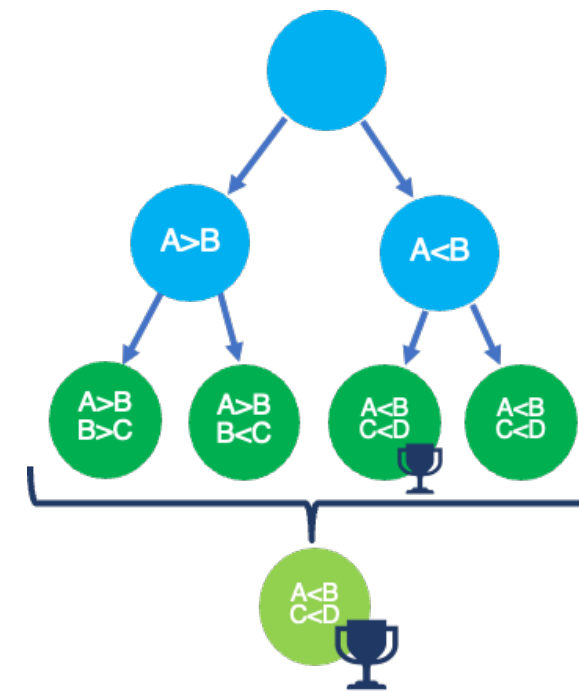
Our three MAPF mechanisms

P-CBS: payment
conflict-based search



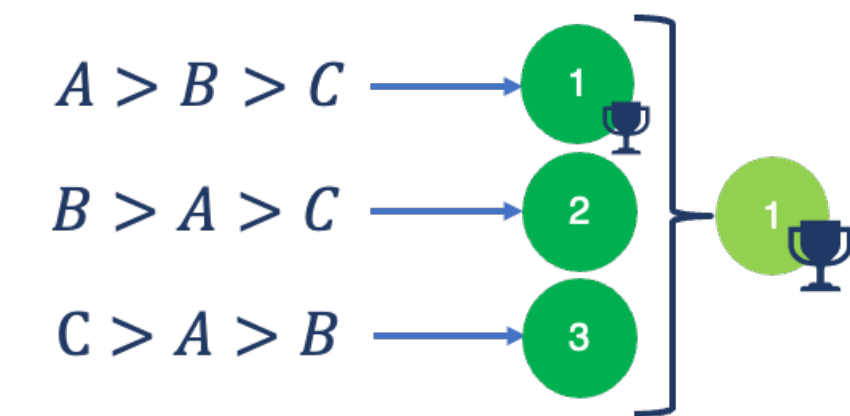
- Optimal MAPF
- VCG payments

E-PBS: exhaustive
priority-based search



- Suboptimal, made strategyproof with MIR
- VCG-based payments "for free"

MC-PP: Monte-Carlo
prioritized planning



- Order agents by priority in m random ways
- Trade off scalability and optimality by choosing m
- VCG-based payments "for free"

Key takeaways

We highlight the **non-cooperative MAPF** domain

- **Standard MAPF** setting, large scale
- **Self-interested** agents

We make **fast MAPF algos strategyproof** with **maximal-in-range (MIR)**:

- Ensure that the allocation is optimal within a **fixed range of outcomes**
 - Add **VCG-based payments** at **no computational cost**
- Experiments on 2D & 3D MAPF benchmark maps validate claims

Future work:

- Apply maximal-in-range (MIR) to other scalable MAPF algos
- Find better ranges to improve suboptimality, e.g. using learning

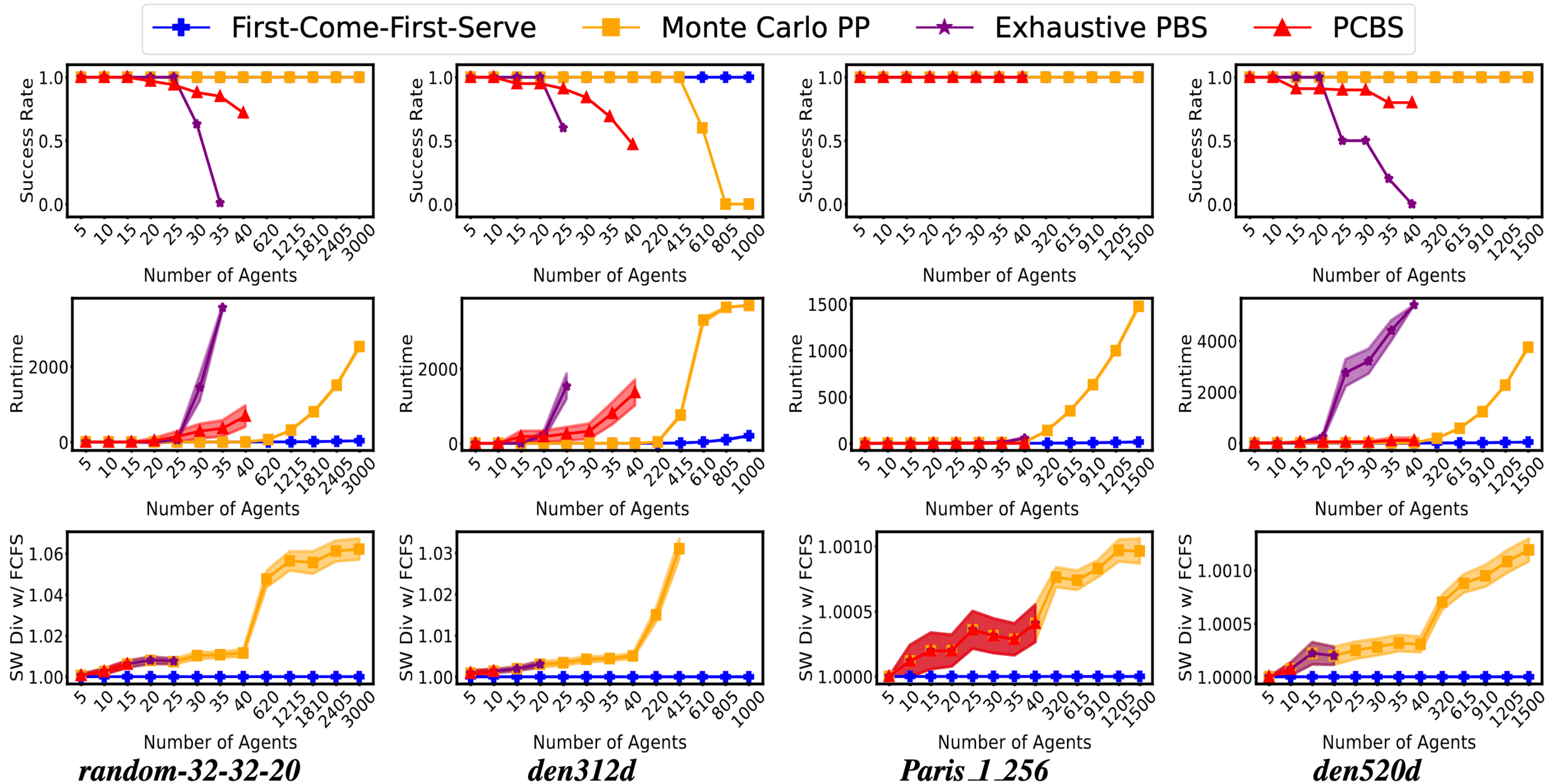
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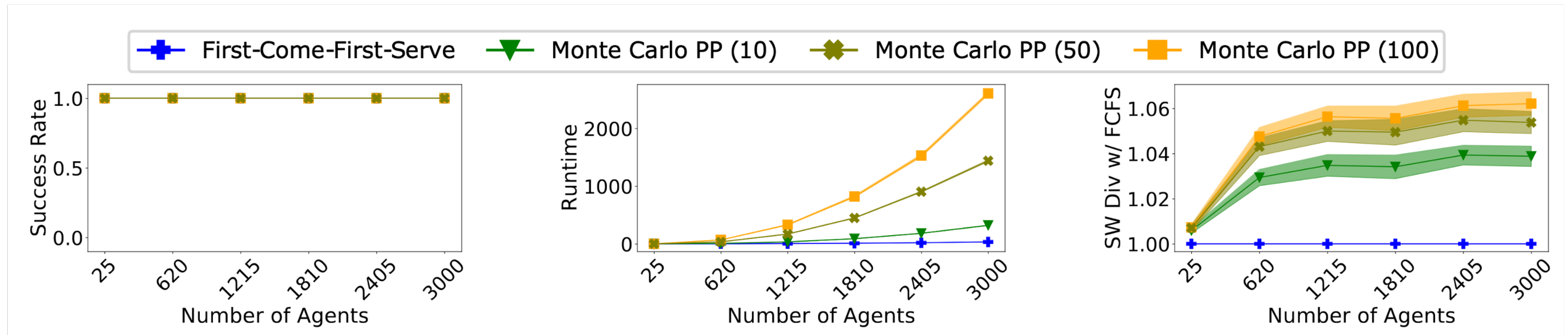
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Our mechanisms on MAPF benchmarks



MCPP: trade off scalability and optimality



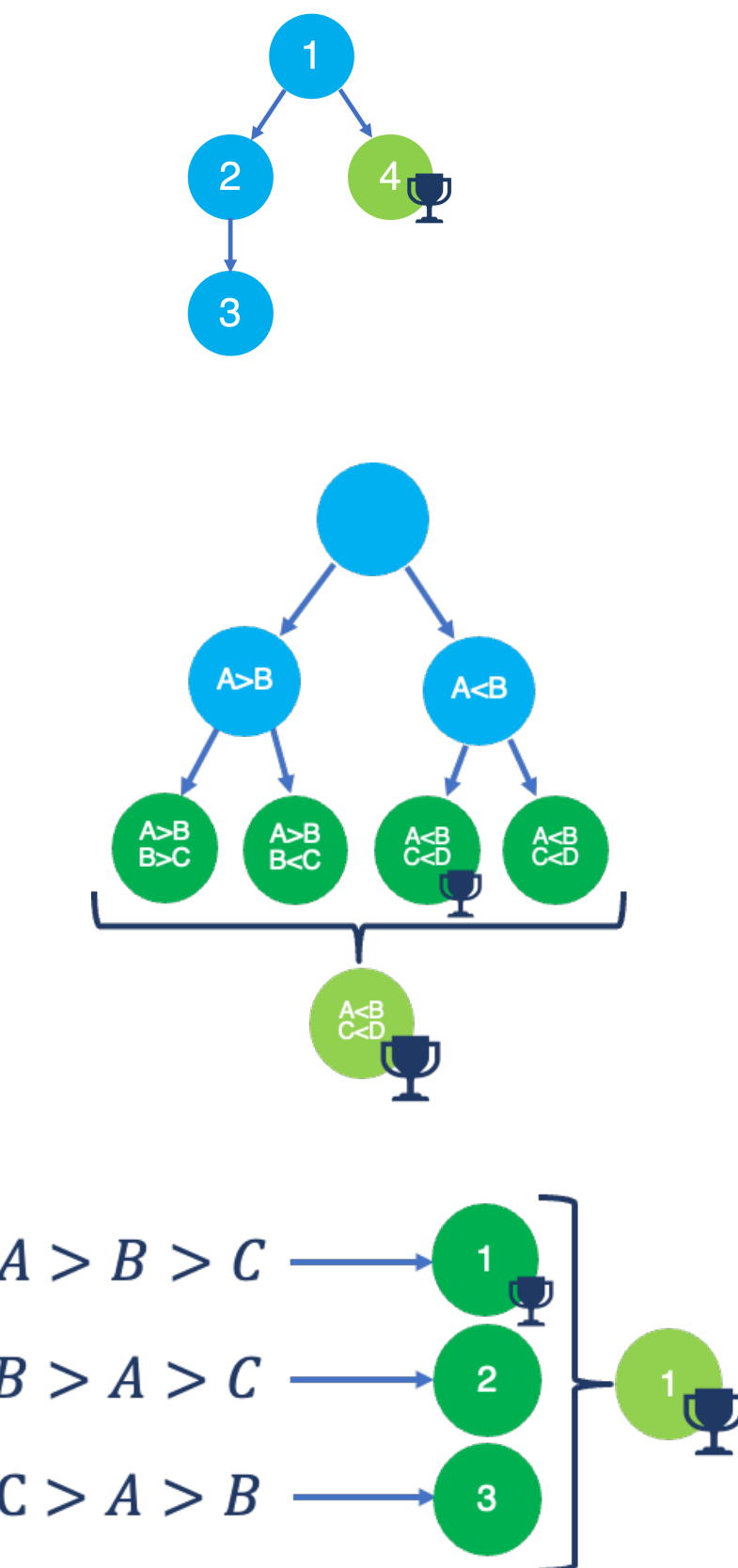
Our three MAPF mechanisms

PCBS: payment conflict-based search

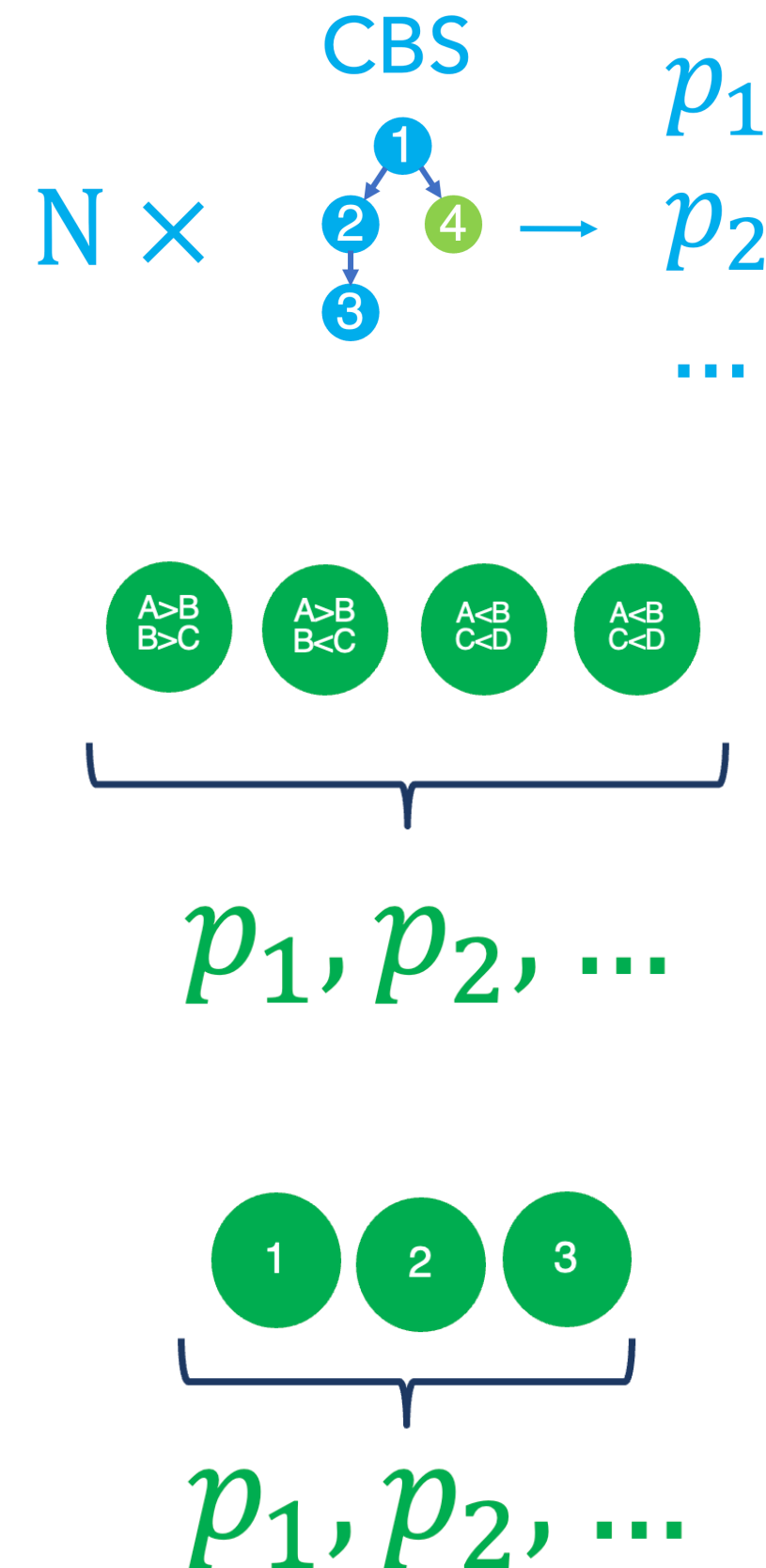
EPBS: exhaustive priority-based search

MCP: Monte-Carlo prioritized planning

Allocation



Payments



Purpose

- Optimal benchmark
- Slow, but parallelizable
- How to make suboptimal MAPF strategy-proof with MIR
- Payments are "free"
- Trade off scalability and optimality via choice of samples
- Parallelizable down to n times A^*

Agent model & optimisation objective

Let $d = (\pi_1, \pi_2, \dots, \pi_N) \in D$ an allocation, i.e. a set of conflict-free paths

MAPF:

- Minimise **sum-of-cost** (or "flowtime"): $\arg \min_{d \in D} \sum_{i=1}^N |\pi_i^d|$
- Minimize **makespan**: $\arg \min_{d \in D} (\max_{i \leq N} |\pi_i^d|)$

Our model: each agent reports a value for arrival v_i and cost per time travelled c_i

- Maximise **social welfare**: $\arg \max_{d \in D} \sum_{i=1}^N v_i - c_i |\pi_i^d|$
- Generalisation of sum-of-cost, identical for $v=0, c=1$